A 25 Gb/s Burst-Mode Receiver for Low Latency Photonic Switch Networks

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Abstract—We report a dc-coupled burst-mode (BM) receiver for optical links in a dynamically reconfigurable network. Through the introduction of interlocking search algorithms, a robust 25 Gb/s BM operation is achieved with 31 ns lock time. At the beginning of the burst, the receiver first performs input dc current offset calibration in 12.5 ns, then achieves phase lock in 18.5 ns, and after that tracks data using a phase interpolator (PI) based bang-bang clock and data recovery (CDR). The sensitivity of the receiver is -10.9 dBm (average power, BER $<10^{-12}$) at 25 Gb/s, tested with a single mode 1550 nm reference optical transmitter. There is no significant sensitivity penalty in the presence of ± 100 ppm frequency offset between the transmitter and the receiver. Measured power efficiency of the receiver at 25 Gb/s is 4.4 pJ/bit. The core of the 32 nm SOI CMOS circuit occupies 200 μ m \times 300 μ m.

Index Terms—Burst-mode (BM), clock and data recovery (CDR), optical receiver, optical switches, silicon photonics, transimpedance amplifier (TIA).

I. INTRODUCTION

R APIDLY reconfigurable photonic switch networks keep the data in the optical domain, bypassing the electro-optical conversion steps and therefore potentially offering significant advantages in latency, bandwidth, and power dissipation [1], [2]. The technical feasibility and scalability of these networks is largely determined by the availability of rapidly reconfigurable, high-port-count optical switches with low insertion loss. Microelectromechanical system (MEMS)-based

Manuscript received May 11, 2015; revised August 04, 2015; accepted September 03, 2015. Date of publication October 09, 2015; date of current version November 24, 2015. This paper was approved by Guest Editor Jaeha Kim. This work was supported in part by the Defense Advanced Research Projects Agency (DARPA) and the Army Research Laboratory (ARL) under contract W911NF-12-2-0051. The views, opinions, and/or findings contained in this article are those of the authors and should not be interpreted as representing the official views or policies of the Department of Defense or the U.S. Government. Approved for Public Release, Distribution Unlimited.

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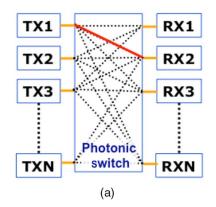
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Digital Object Identifier 10.1109/JSSC.2015.2478837

optical switches scalable to over 1000 ports have been demonstrated [3], [4], but their millisecond-scale reconfiguration times greatly exceed the target message latencies for bigdata applications [2]. Liquid crystal on silicon (LCOS)-based switches [5] are scalable to $\sim \! 100$ ports and offer a significant improvement in reconfiguration speed, but they are still limited to microsecond-scale operation times. The fastest optical switches demonstrated to date are based on silicon photonics [6]. Although these switches have a limited scalability ($\sim \! 10^{\circ}$ s of ports), they were shown to operate on nanosecond-scale times. Full utilization of the speed of these switches, however, requires nanosecond-scale burst-mode (BM) transceivers, performing a function similar to the one required in passive optical networks (PON), but on a much faster scale.

In this work, we report a BM receiver for optical links in a dynamically reconfigurable network. To support the extremely high bandwidth of the envisioned optically switched network (10 Tbytes per second per node [2]), and to provide a high degree of flexibility, the receiver data rate target has to cover the range of 15–25 Gb/s. System-level simulation of the network [2] showed that receiver locking time of at most 100 ns would be required, and locking times under 50 ns would be highly desired, with faster locking times resulting in further improvement of the overall network performance. The main challenge for the receiver design stems from the fact that after the switching event the receiver has to lock to the signal that is not only originating from a different transmitter, but also taking a different path and passing through a different set of switches in the network, as illustrated in Fig. 1. The optical signals originating from different transmitters have different extinction ratios and accumulate different levels of loss as they travel through the network. To accommodate this variability, the receiver has to support at least 10 dB of dynamic range in the input optical power. The corresponding offset compensation calibration has to be performed on a nanosecond scale, within the sub-50 ns total locking budget. In addition, the receiver has to support uncoded messages of arbitrary length [2], requiring the frontend to be dc-coupled. Consequently, the dc offset compensation setting of the receiver found in BM has to be applied indefinitely, for the duration of the message, resulting in a "no low frequency cutoff" requirement. In order to enable scalability to very large systems, the receiver has to support ± 100 ppm frequency offsets between the transmitter and the receiver clocks, which can be potentially generated from independent crystal reference sources. The overall receiver architecture also has to provide a path for support of much larger frequency offsets,



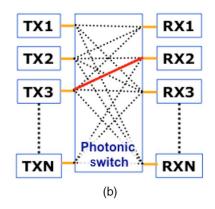


Fig. 1. Example of two different states of a photonic switch network, connecting receiver RX2 to (a) transmitter TX1 or (b) transmitter TX3.

resulting, e.g., from spread spectrum clocking requirement, typical for large computing systems. Finally, the receiver has to maintain an aggressive link power efficiency target (under 5 pJ/bit). The overall switching network was envisioned as utilizing semiconductor optical amplifiers (SOAs) to compensate for switch losses [2], so the overall power dissipation of the system is a function of the receiver sensitivity. The sensitivity goal was set to -10 dBm average optical power, or better.

The majority of BM optical receivers reported in the literature address PON applications (for a recent review, see [7]). While there are many similarities between the goals of our work and PON receiver standards, there are also a number of key differences. Similarly to PON, at the beginning of the burst event, the receiver in a photonic switch network has to first calibrate the transimpedance amplifier (TIA) and then engage the BM clock and data recovery (BM-CDR) circuit to recover and track the phase of the incoming signal. The differences originate from the fact that, first, there are no PON standards covering our target data rate in the range of 15–25 Gb/s and, second, PON places a significant stress on receiver sensitivity, while locking time requirements are relatively relaxed. The most demanding specification of the 10G-EPON standard, e.g., is the combination of -28 dBm sensitivity with a 22 dB dynamic range of the optical signal, which far exceeds the performance goals of the receiver in a photonic switch network, as outlined above. At the same time, the lock time target of the 10G-EPON standard is set at 800 ns for the BM front-end settling plus another 400 ns for the CDR lock, dramatically slower than the target lock time for the BM receiver in a photonic switch network. Even the fastest reported lock time of 75 ns [8] for the BM front-end still does not reach our goal of sub-50 ns total time (including the CDR lock time). For comparison purposes, we can assume that locking speed scales linearly with data rate. In that case, 75 ns at 10 Gb/s will become 30 ns at 25 Gb/s, which is closer to our target but still does not quite reach the desired level of performance. In this work, we report a BM TIA calibration scheme that completes in 12.5 ns. The TIA calibration state machine employs a successive approximation search algorithm and holds the identified offset values in the digital domain, enabling full compatibility with the requirements for dc-coupling and arbitrary message length.

The BM-CDR problem can be solved in a number of different ways. In PON applications, a popular approach is based on gated voltage-controlled oscillator (G-VCO) architecture [9]–[11]. While this approach allows fast lock times, it has very poor jitter rejection characteristics. For a photonic switch network, we were looking for a CDR that would support a wide range of data rates, from 15 to 25 Gb/s, and that would require a G-VCO with a very wide tuning range. As is well known, in VCOs, there is a direct tradeoff between tuning range and phase noise, further degrading the expected jitter performance of the G-VCO-based CDR. Finally, G-VCO architecture requires a dedicated VCO per channel, significantly increasing power dissipation and potentially leading to cross-talk issues in a multichannel receiver. The BM-CDR architecture based on phase picking [12] instantaneously identifies an acceptable sampling phase from a small fixed set of available clock phases. While the selected phase allows correct sampling of the data, the jitter properties of the recovered clock remain very poor as it experiences large random phase jumps. A fast CDR locking mechanism was proposed in [13], based on the idea of generating correct sampling phase by interpolating between the quadrature clocks. The main concern with this approach requiring a full rate clock, full rate sampling, and an analog full rate phase interpolator (PI) is its scalability to 25 Gb/s. A more direct, mostly digital solution of the BM-CDR problem is offered by the oversampling architecture [14]. The oversampling CDR effectively uses a flash time-to-digital converter (TDC) to identify the positions of the data transitions and the optimum phases for data sampling. To ensure a smooth hand-off of the identified data transition phases to the regular bang-bang CDR (BB-CDR) for tracking and to avoid the long locking transient typical of bang-bang control loops, the spacing between the phases has to be very small. The hand-off should place the CDR where it would normally be when in lock, i.e., within the region where the bang-bang phase detector is effectively linearized by the jitter. As a result, the data edge position has to be located with accuracy on the order of 1–2 ps (or less, depending on the jitter content of the incoming data), requiring a very high oversampling ratio. While this is certainly a very robust, simple, and scalable architecture, the need for a large number of sampling phases dramatically increases the power dissipation of the overall system. In this work, we propose a new digital BM PI-based CDR architecture that uses a small number of clock phases ($3 \times$ oversampling) and a successive approximation algorithm to search for the data transition. Compared to the "flash TDC" oversampling BM-CDR, the proposed "SAR-TDC" BM-CDR architecture trades off a small increase in lock

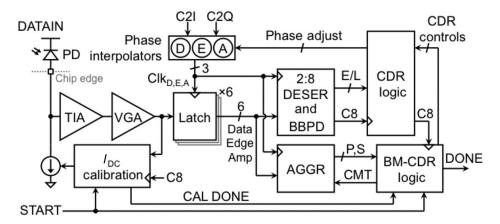


Fig. 2. Burst-mode receiver architecture.

time for a significant reduction of power dissipation and simplification of the overall design, while still keeping all the benefits of the predominantly digital nature of the approach. The BM-CDR search algorithm starts upon completion of TIA calibration, achieves lock in 18.5 ns, and hands off the control to the regular bang-bang PI CDR for tracking. The overall locking time of the 25 Gb/s receiver is 31 ns, with 4.4 pJ/b efficiency and -10.9 dBm average optical power sensitivity. The receiver operates without any significant sensitivity penalty in the presence of ± 100 ppm frequency offset between the transmitter and the receiver. The architecture of the CDR allows straightforward modifications (addition of low speed digital integrators) to support much larger frequency offsets and spread spectrum clocking.

II. BM RECEIVER ARCHITECTURE

The block diagram of the receiver is shown in Fig. 2. The photocurrent from the dc-coupled photodiode (PD) is converted into voltage by the TIA. This voltage is then amplified by a variable gain amplifier (VGA) and sampled by three groups of latches (Edge, Data, and Amp). The latches are clocked by three half-rate clock phases (Clk_E, Clk_D, and Clk_A), generated by three independent PIs (PI_{E,D,A}) from in-phase (C2I) and quadrature (C2Q) clocks (generated on-chip using a static frequency divider). The Edge, Data, and Amp samples are deserialized further (2:8 DESER) and processed by the bang-bang clock and data recovery (CDR) circuit. The structure of the PI, latches, deserializer, and the BB-CDR logic blocks is similar to those described in [15]. The receiver has three control loops operating in sequence at the beginning of each burst, triggered by the rising edge of the START signal, as illustrated in Fig. 3. In the envisioned network, the START signal (Step 1 in Fig. 3) is provided by a central controller, which also reconfigures the transmitter and the switch. The TIA calibration engine sets the dc offset current I_{DC} at the input of the TIA (Step 2 in Fig. 3), and, upon completion, asserts the CAL DONE signal, initiating the BM-CDR control loop. After that, the BM-CDR locks to the incoming data by taking over the PI controls, processing the Edge, Data, and Amp samples in the aggregator (AGGR), and directly adjusting the phases of all three PIs (Step 3 in Fig. 3). Finally, upon completion, the BM-CDR logic asserts a DONE

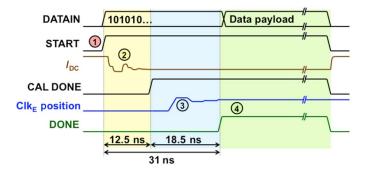


Fig. 3. BM timing diagram: 1) start of new burst, 2) IDC calibration converges on decision threshold, 3) BM-CDR aligns Clk_E to DATAIN edge, and 4) BB-CDR takes over, Clk_D samples valid data.

signal and returns PI controls to the regular CDR loop (Step 4 in Fig. 3). The end of the burst is marked by the falling edge of the START signal, which resets all state machines.

A. Analog Front-End

A detailed schematic of the analog front-end of the receiver is shown in Fig. 4. The TIA is realized as a CMOS inverter with resistive feedback, as shown in Fig. 5(a). This TIA topology is chosen because of its excellent gain and noise at low power consumption [16]. The TIA power supply (VDD_TIA) is separate from the other power supplies to allow independent control of TIA biasing and to provide power supply noise isolation from other supplies which have digital switching activity. The TIA feedback resistor $R_{\rm FB}$ is controlled by a 4 bit thermometercoded signal and can be varied from 203 to 299 Ω at nominal process. The range provides sufficient control such that the resistor value can be adjusted to $250 \pm 13 \Omega$ across process, voltage, and temperature (PVT) variations. Postlayout extraction simulation (nominal process, 25 °C temperature, $+1.5\sigma$ wire parasitic capacitance, VDD_TIA = 1 V, 80 fF PD capacitance, and 200 pH wirebond inductance) indicates $46~\Omega$ input impedance, 210 Ω transimpedance gain, 18.4 GHz bandwidth, and 2.7 mW power dissipation of the TIA.

A replica TIA provides the dc voltage comparison point for the differential VGAs. The replica TIA also provides commonmode power supply rejection at low frequencies while most of

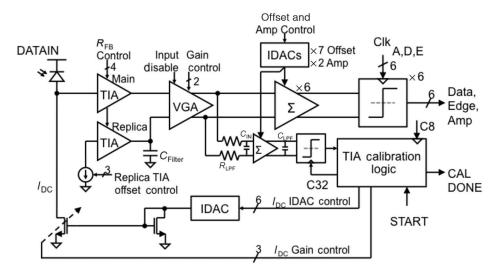


Fig. 4. Analog front-end schematic.

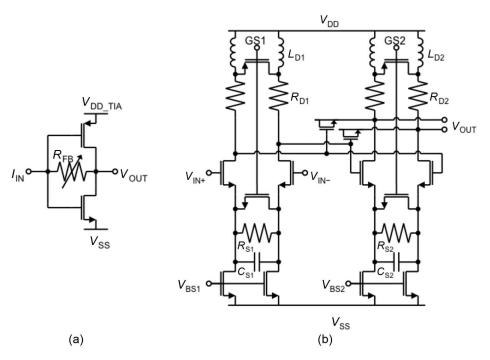


Fig. 5. Schematic of (a) TIA and (b) VGA.

 $TABLE\ I \\ POSTLAYOUT\ EXTRACTED\ SIMULATION\ OF\ VGA\ GAIN\ AND\ BANDWIDTH$

Gain setting	0.5dB (1.06×)	5.5dB (1.89×)	10.7dB (3.43×)
Bandwidth (GHz)	24.1	23.9	22.1

the thermal noise of the replica TIA above 4 GHz is filtered out by the capacitor $C_{\rm Filter}$ (Fig. 4). The simulated input-referred current mismatch between the main and replica TIAs is normally distributed with $\sigma=40~\mu{\rm A}$. The dc component of the input current $I_{\rm DC}$, automatically selected at the start of each burst (detailed discussion below), is also used to compensate for the offset of the main TIA with respect to the replica TIA. Should $I_{\rm DC}$ approach 0, a coarse 3 bit replica TIA offset control can apply up to $140~\mu{\rm A}$ (over 3σ of mismatch) to the input of the replica TIA, allowing $I_{\rm DC}$ to be brought back into range.

The VGA [Fig. 5(b)] has two bits of gain control and an input disable feature to allow latch offset compensation. The gain of each of the two stages of the VGA can be digitally switched between $5.4~(1.85\times)$ and $0.2~dB~(1.03\times)$ using control signals GS1 and GS2. The variable gain is used to prevent limiting behavior in the VGA. Limiting amplifier action would result in voltage noise to jitter conversion and conversion of dc offset into duty cycle distortion. In the 5.4~dB gain state, the inductors are used for bandwidth extension. In the 0.2~dB

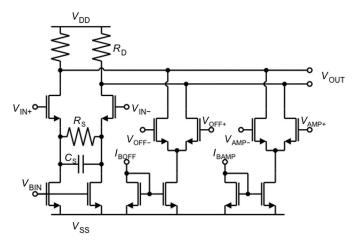


Fig. 6. Schematic of the summer.

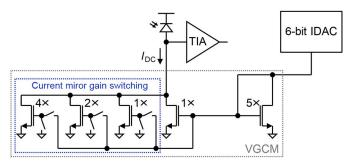


Fig. 7. VGCM schematic.

gain state, the inductors are shorted out differentially and RC source degeneration is used to extend the bandwidth. By using different bandwidth extension techniques, we can optimize for the maximum possible bandwidth extension in each gain state. The VGA bandwidth is also extended by adding negative Miller capacitors to the second stage, as shown in Fig. 5(b). To disable the input, the biasing circuit (not shown) sets $V_{\rm BS1}=0~{\rm V}$, turning OFF the first stage of the VGA and causing the second stage to output a differential 0 at V_{OUT} . Note that in this mode ($V_{BS1} = 0 \text{ V}$), the output of the first stage of the VGA is at VDD, which suppresses the offsets in the second stage of the VGA. At the same time, the common mode of the output of the second VGA stage will be close to the common mode of the signal in normal operation. As a result, setting $V_{\rm BS1}=0~{\rm V}$ provides a differential zero at a common mode close to the correct value, enabling offset correction for each latch with a corresponding summer individually. Postlayout extraction simulations at nominal process, 25 °C temperature, and VDD = 1 V with $+1.5\sigma$ wire parasitic capacitance result in 11.8 mW power, with gain and bandwidth at different settings listed in Table I. The simulated total input-referred current noise of the TIA and VGA is 2.59 µArms. This corresponds to the expected theoretical sensitivity of -14.4 dBm (average power) at BER = 10^{-12} , assuming 0.5 A/W responsivity of the PD.

Following the VGA are the summers (Fig. 6), which use current summation into a resistor load. The summers add a controlled amount of dc offset to the input signal, enabling device mismatch compensation and input amplitude monitoring. A set of nine IDACs provide digitally controlled currents

to the summers, seven 5 bit IDACs for offset compensation and two 6 bit IDACs for amplitude monitoring. Six summers are driving the high-speed path latches (Edge, Data, and Amp), while the seventh summer is used in the TIA calibration feedback path latch (as shown in Fig. 4). The 5 bit IDACs are used in all seven summers, while the 6 bit IDACs are only used in the two summers in the Amp path. The latches use a modified StrongARM sense amplifier topology [17], [18]. The summer consists of an input differential pair with RC degeneration to increase the input-to-output bandwidth and two current mirrors with polarity-steering switches to inject offset and amplitude currents into the output nodes. The offset and amplitude IDACs are connected to the I_{BOFF} and I_{BAMP} inputs, respectively. The input, offset, and amplitude currents are summed at the output nodes and converted to a voltage by the load resistors to produce the final output voltage. The load resistors (R_D in Fig. 6) are set to $315~\Omega$ to provide sufficiently fast settling for kickback from the latch. Postlayout extraction simulation at nominal process, 25 °C temperature, and VDD = 1 V with $+1.5\sigma$ wire parasitic capacitance results in 2.1 dB (1.27×) gain, 31.7 GHz bandwidth (with latch in reset), and power ranging from 1.6 to 3.3 mW, depending on IDAC settings. The simulation with the clock running has shown that latch common-mode kickback settles to within 30 mV before the next sampling, while differential kickback settles to less than 4 mV error across PVT corners.

B. BM TIA Calibration Engine

The TIA calibration logic removes the dc component of the photocurrent (I_{DC}) at the input of the TIA by first selecting the ratio in a variable gain current mirror (VGCM) and then using a binary search algorithm to find the optimal 6 bit IDAC setting. The proposed scheme is designed to compensate for dc currents up to 1 mA (exceeding the 10 dB dynamic range goal of the receiver), while keeping the resolution at 2%–4% of the peak-to-peak ac current amplitude, to minimize the sensitivity penalty due to the residual offset. For example, at 50 µA peak-to-peak ac current, the 4% (2 μA) resolution would result in sensitivity penalty of only 0.087 dB (maximum offset of $\pm 1~\mu$ A). In the envisioned network, the transmitter extinction ratios will vary in a wide range, but they should always be better than 5 dB. As a result, we expect the peak-to-peak ac current amplitude to scale with the value of I_{DC} . This consideration allows a 6 bit IDAC with 4 setting (3 bit thermometer-coded) gain control to achieve both maximum 1 mA range and minimum 2 µA resolution. Note that if we required a constant resolution of 2 μA throughout the 1 mA range, we would need a significantly larger 9 bit IDAC. We use a standard currentsteering IDAC architecture, which is both fast and low glitch. The IDAC current source transistors are sized to minimize noise impact. The 6 bit IDAC least significant bit (LSB) size is 10 μA and the maximum output current is 630 µA.

The VGCM schematic is shown in Fig. 7. The current mirror adjusts gain by changing the number of current sources on the output side of the mirror; this helps maintain large gate-to-source voltage across the diode-connected transistor on the input side of the mirror, improving VGCM's speed and

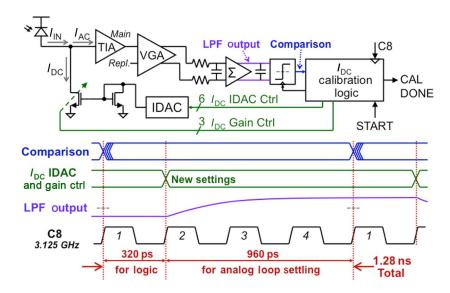


Fig. 8. $I_{\rm DC}$ calibration timing diagram.

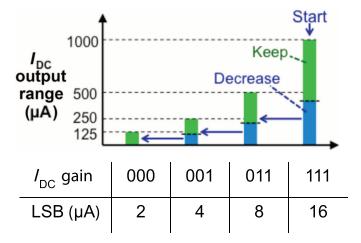


Fig. 9. I_{DC} calibration algorithm.

noise performance. The input-to-output current mirror ratio can be set to 5:1, 5:2, 5:4, or 5:8. Together with the 10 μ A 6 bit IDAC LSB, the VGCM provides a minimum $I_{\rm DC}$ LSB of 2 μ A at 5:1 gain and a maximum $I_{\rm DC}$ output of 1.008 mA at 5:8 gain, thus satisfying the $I_{\rm DC}$ range and resolution requirements.

To minimize locking time, both the TIA calibration engine and BM-CDR require a " $1010\ldots$ " pattern preamble from the rising edge of the START signal until DONE is asserted. To determine the dc level of the input signal, a low-pass filter (LPF) is used to filter out the " $101010\ldots$ " preamble. The LPF has two poles, one created by $R_{\rm LPF}$ and the summer's input capacitance, the other created by the summer's output resistance and $C_{\rm LPF}$. The two pole frequencies are in the 1–2 GHz range to filter the 12.5 GHz " $101010\ldots$ " signal while still responding rapidly to changes in the $I_{\rm DC}$ value. Across the PVT corners, the simulated filtering at 12.5 GHz exceeds 30 dB. The simulated 2% settling time for changes in $I_{\rm DC}$ is less than 0.98 ns, across PVT variations. The output of the latch following the LPF is used as the input to the TIA calibration engine to determine if $I_{\rm DC}$ should be increased or decreased.

The timing diagram of the I_{DC} calibration engine is shown in Fig. 8. The state machine runs on the C8 clock (1/8th of the data rate, or 3.125 GHz at 25 Gb/s), so the speed of the locking time scales linearly with data rate. The calibration path latch is clocked at the beginning of C8 cycle 1, producing the "Comparison" output. Once the "Comparison" bit is valid, the $I_{\rm DC}$ calibration logic computes the new settings for the IDAC and gain, which are made available to the IDAC and VGCM at the beginning of cycle 2 in Fig. 8. The analog components then have until the end of cycle 4 to settle, at which time the latch is clocked again. At a C8 clock period of 320 ps, a single step of the I_{DC} calibration logic takes $4 \times 320 \text{ ps} = 1.28 \text{ ns. A maxi-}$ mum of three steps are needed to determine the gain setting and six steps are required to determine the IDAC setting, resulting in a total of nine steps, or 11.52 ns, to set I_{DC} . To synchronize the START signal with the I_{DC} calibration logic, we use triple latching to resolve any potential metastability issues, which requires three more C8 cycles and results in a total time of 12.5 ns from START signal going high to I_{DC} being determined and CAL DONE going high.

The I_{DC} gain search algorithm flowchart is illustrated in Fig. 9. The VGCM gain is initially set to the highest value (111), and the IDAC is set to 41% of maximum to provide range overlap in case of gain errors. At each VGCM gain setting, the dc value of the input signal is compared to 41% of the maximum I_{DC} at this gain setting. If the dc value is less than 41% of the maximum, the gain is decreased to the next value and the search continues; otherwise, the gain setting is found. The gain search also stops when it reaches the lowest setting of 000. The gain search takes, at most, three steps to complete. Once the gain setting is determined, a six step binary search algorithm is used to identify the 6 bit IDAC value. The BM TIA calibration engine completes the search for VGCM gain and IDAC settings in 12.5 ns (at 25 Gb/s) and hands over the control to the BM-CDR state machine by asserting the CAL DONE signal. From a system point of view, it is important to point out that the resulting setting of I_{DC} is available in digital form, which can be stored and reapplied later when the same

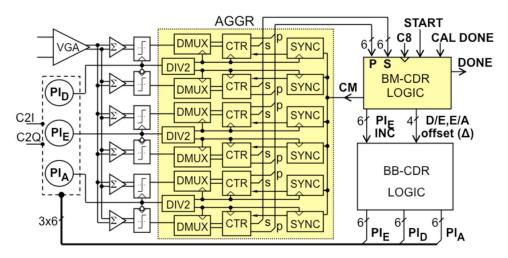


Fig. 10. BM-CDR block diagram.

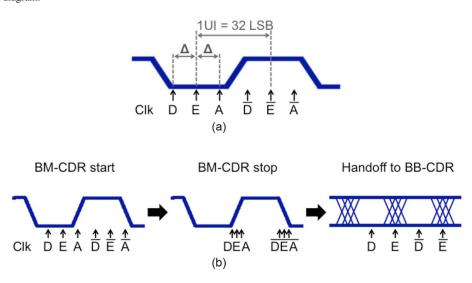


Fig. 11. (a) Definition of one LSB of the PI and of the D,A phase offset Δ . (b) PI positions: at the start of the BM-CDR search, upon BM-CDR completion and at the handoff to BB-CDR.

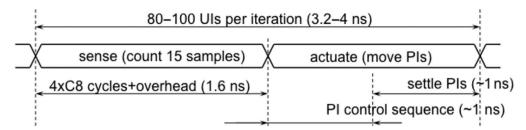


Fig. 12. BM-CDR iteration cycle and timing.

transmitter and the same switch configuration are used in order to significantly reduce or even completely eliminate the BM TIA calibration time.

C. Burst-Mode CDR

The BM-CDR block diagram is shown in Fig. 10. The BM-CDR state machine controls both the AGGR and the PIs, overriding the regular BB-CDR logic. In this mode, the BB-CDR macro is used as a passive interface to apply the control signals issued by the BM-CDR to the PIs. The BM-CDR

finds the data transition edge using a successive approximation algorithm by estimating the position of the edge based on accumulated samples and setting a progressively narrower window for sampling data on both sides of the transition. The search for the data transition edge starts with PIs PI_D and PI_A offset from the PI_E by $\Delta=11$ LSB each (32 LSBs equal 1 bit time, or 1 UI). The definitions of PI LSB and offset Δ are illustrated in Fig. 11(a). This initialization choice spreads the 6 half-rate sampling points approximately evenly across 2 UI [Fig. 11(b)]. The final BM-CDR goal position of the PIs is

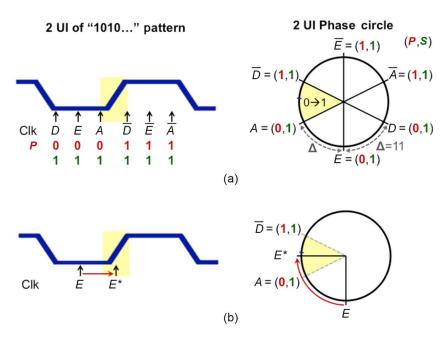


Fig. 13. BM-CDR (a) locates the $0 \to 1$ transition sector and (b) places PI_E inside the identified $0 \to 1$ transition sector.

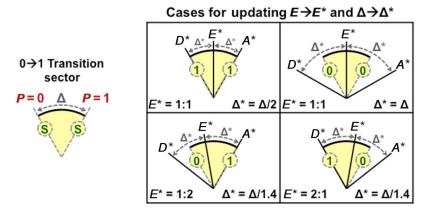


Fig. 14. The four possible values of the S-bits on the edges of the $0 \to 1$ transition sector determine the new PI_E position E^* and the new value of $\Delta(\Delta^*)$.

TABLE II BM-CDR CONVERGENCE RULES

Δ	Δ/3	Δ/2	2∆/3
11	3	5	7
8	2	4	5
6	2	3	4
4	1	2	3
3	1	1	2
2	1	1	1

 ${\rm PI_{A,D}}$ offsets Δ from ${\rm PI_E}$ in units of LSB and the possible new ${\rm PI_E}$ positions $(\Delta/3,~\Delta/2,~2\Delta/3)$, selected based on the values of S-bits

also shown in Fig. 11(b). The function of the AGGR is to integrate 15 consecutive binary samples from all six latches and to present the result as a set of six 1 bit pairs reflecting polarity (P) and saturation (S) information. If the sum of the 15 samples is 7 or less, then P is set to 0; otherwise, it is set to 1. S is set to 1 if the sum exactly equals 0 or 15, and it is 0 otherwise. The

BM-CDR iteration cycle and timing are illustrated in Fig. 12. The BM-CDR first issues a command to integrate 15 samples ("sense" in Fig. 12). Then, based on the result, it moves PI_E to an estimated position of the data transition, and places the other two PIs as left and right guard-bands ("actuate" in Fig. 12). The iteration cycle is split approximately 50/50 between sensing (counting samples) and actuation (moving PIs to a new position).

The BM-CDR first finds a pair of phases corresponding to a $0 \to 1$ transition in polarity bit P as illustrated in Fig. 13(a). In this example, the $0 \to 1$ transition sector is located between the phases A and \bar{D} . Then, the algorithm estimates the location of the new position of PI_E and places the E phase between those two phases as shown in Fig. 13(b). The exact placement of the E phase depends on the values of the two saturation bits S on the edges of the $0 \to 1$ transition sector (1,1 in the example shown in Fig. 13). The values of the S-bits are also used to decide on the new value of Δ for the next iteration. The BM-CDR convergence rules are illustrated in Fig. 14 and summarized in Table II. If the pair of saturation bits S on the sides of

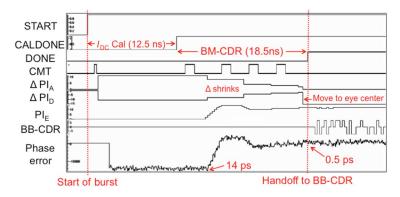


Fig. 15. Verilog simulation of phase locking dynamics at 25 Gb/s.

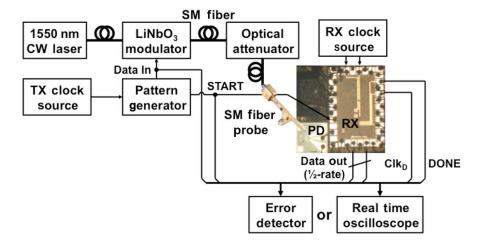


Fig. 16. Test setup.

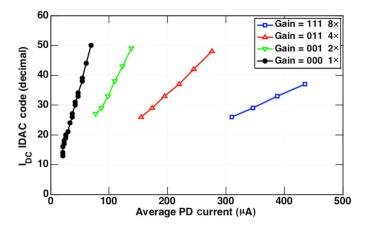


Fig. 17. Measured $I_{\rm DC}$ calibration values with different input optical power levels.

the $0 \to 1$ transition is either 0,0 or 1,1, the new edge position is set symmetrically 1:1 at midpoint (half of the current phase separation Δ : the " $\Delta/2$ " column in Table II). If that pair is 0,1 or 1,0, the edge position is set closer to the S=0 point, at 1:2 or 2:1 ratio accordingly (1/3 or 2/3 of the current phase separation Δ : the " $\Delta/3$ " and " $2\Delta/3$ " columns in Table II). The starting position is $\Delta=11$ LSB (top line in Table II). Depending on the values of the S-bits, Δ is either left unchanged (if the sum

of S-bits is 0), or it is reduced by $1.4\times$ (moving down one line in the table, if the sum of S-bits is 1), or it is reduced by $2\times$ (moving down two lines in the table, if the sum of S-bits is 2). The algorithm stops when Δ reaches two LSB (bottom line in Table II), or upon timeout. The BM-CDR completes the search in 18.5 ns (at 25 Gb/s), asserts the DONE bit and hands over the control of the PIs to the regular BB-CDR, keeping PI_E in place and moving PI_D into the middle of the eye ($\Delta = 16$ LSB), as illustrated in Fig. 11(b).

The BM-CDR search algorithm was verified in Verilog simulation (Fig. 15). As the BM-CDR search progresses, the digital offsets Δ of PIs PI_{A,D} relative to PI_E converge to a progressively smaller value. The bottom trace ("phase error" in Fig. 15) is defined as the absolute Clk_E to DATAIN edge-to-edge distance (in picoseconds). The dynamic of the "phase error" corresponds to the dynamic of the digital control of the PIE PI in Fig. 15. The BB-CDR (enabled after the DONE signal) shows balanced early/late signals and keeps PIs in the same position, confirming the correctness of the phase lock. The Verilog simulation was run for all possible initial values of the "phase error" with 1 ps step, at different values of random jitter on DATAIN and the receiver clock. The simulations have shown that in all scenarios the BM-CDR behaved correctly and found the target CDR state in under 19 ns. Correct BM-CDR dynamics were also verified in the presence of ± 100 ppm frequency offsets between DATAIN and the receiver clock. But

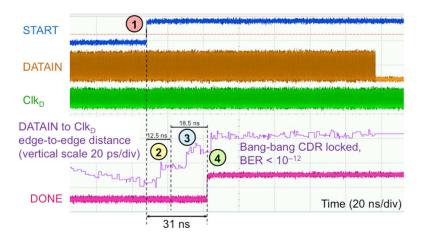


Fig. 18. Measured 25 Gb/s burst waveforms showing phase locking dynamics: 1) start of burst, 2) I_{DC} calibration and PI reset, 3) BM-CDR acquisition, and 4) Clk_D at eye center, BB-CDR locked.

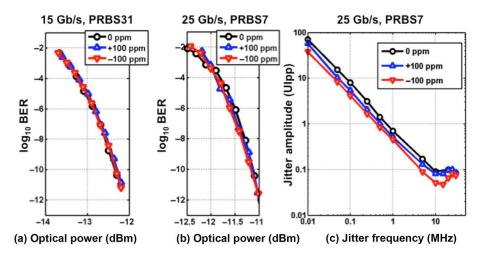


Fig. 19. Measured sensitivity at (a) 15 Gb/s and (b) 25 Gb/s, with and without frequency offsets. (c) Measured jitter tolerance at 25 Gb/s.

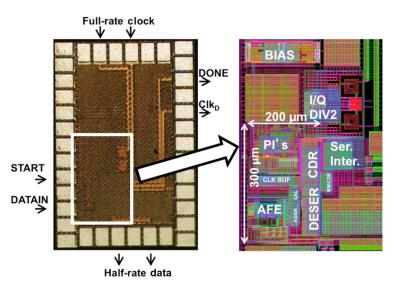


Fig. 20. Die photo with layout detail.

it should be noted that the additional phase difference arising from ± 100 ppm frequency offsets and accumulated while BM-CDR is running is less than 2 ps, easily absorbed by the search algorithm.

III. EXPERIMENTAL RESULTS

The 32 nm SOI CMOS receiver chip was wirebonded to a commercial PD (40 GHz bandwidth, $12 \mu m$ diameter, $100 \, \text{fF}$ capacitance, and $0.5 \, \text{A/W}$ responsivity) and tested with

Technology	IBM 32 nm SOI CMOS	
Maximum data rate	25 Gb/s	
Area	0.06 mm ²	
Supply voltages	$V_{\rm DD,TIA}$	1.00 V (4 mW)
(with power consumption,	V _{DD,A}	0.94 V (80 mW)
PI _A ON)	V _{DD,D}	1.00 V (25 mW)
Total power consumption	PI _A ON	109 mW
	PI _A OFF	100 mW
Power efficiency at 25Gb/s	PI _A ON	4.4 pJ/bit
	PI _A OFF	4.0 pJ/bit
Sensitivity at BER<10 ⁻¹² , PRBS7	15 Gb/s	-12.0 dBm
(avg. optical power)	20 Gb/s	−11.8 dBm
	25 Gb/s	−10.9 dBm
Lock time at 25 Gb/s	31 ns	

TABLE III
BM RECEIVER PERFORMANCE SUMMARY

a reference transmitter consisting of a 1550 nm laser and a LiNbO₃ modulator, followed by an optical attenuator. The test setup (Fig. 16) featured independent clock sources for the transmitter and the receiver and enabled both BM (I_{DC} calibration, BM-CDR dynamics) and continuous mode (BER, BB-CDR tracking) characterization of the receiver. The measured performance of the $I_{\rm DC}$ calibration engine is shown in Fig. 17. In this experiment, the optical attenuator was stepped from 0 to 13 dB in 0.5 dB steps. The 3 bit gain and 6 bit IDAC values were selected by the I_{DC} calibration engine, and read out using an on-chip serial interface. For a low average optical power at the input corresponding to the lowest setting of the gain control vector, repeated burst measurements resulted in only one LSB variation of the IDAC value. Correct performance of the $I_{\rm DC}$ calibration engine was verified at data rates from 15 to 25 Gb/s by comparing the automatically selected IDAC and gain settings to those manually optimized for best possible BER at a given level of input optical power. The automatically selected values were found to always be within two LSB of the optimum. Fig. 18 shows the measured phase locking dynamics of the BM-CDR at 25 Gb/s. The time between the rising edges of DATAIN (pattern generator signal driving the modulator) and Clk_D was measured by the real-time oscilloscope and plotted together with voltage waveforms. The three jumps in the Clk_D phase in Fig. 18 correspond to 1) PI reset after the START signal; 2) initial rapid acquisition by BM-CDR; and 3) the final placement of Clk_D in the middle of the eye. The DONE signal is asserted 31 ns after the START and after that the relative phase of Clk_D shows no significant change during bang-bang operation. To study the effect of different values of the initial phase between incoming data and local clock, BM operation was verified a number of times in the presence of a small frequency offset between the transmitter and the receiver. The BB-CDR was consistently found to be in lock after the DONE signal was asserted.

The sensitivity of the receiver shown in Fig. 19(a) and (b) was measured to be -10.9 dBm (average power) at 25 Gb/s (BER $< 10^{-12}$, PRBS7). This is about 3.5 dB worse than the

expected theoretical value of -14.4 dBm based on postlayout extracted simulations of the analog front-end. At 15, 20, and 25 Gb/s, there was no significant sensitivity penalty in the presence of ± 100 ppm frequency offset between the transmitter and the receiver. At 15 Gb/s, there was no sensitivity penalty going from PRBS7 to PRBS31, at 20 Gb/s the penalty was 0.6 dB and at 25 Gb/s the PRBS31 penalty was 4.6 dB. The PRBS31 sensitivity penalty is attributed to the imperfections in packaging of the decoupling capacitor for the PD power supply. The CDR jitter tolerance curve in Fig. 19(c) was measured at 25 Gb/s (PRBS7, BER $\sim 4 \times 10^{-12}$) with and without ± 100 ppm frequency offset. Measured power efficiency at 25 Gb/s is 4.4 pJ/bit. In continuous CDR mode, the PIA can be powered down and the efficiency improves to 4 pJ/bit at 25 Gb/s. The die photo annotated with layout detail is shown in Fig. 20. The core of the receiver occupies 200 μ m \times 300 μ m. The total number of transistors in the design was 63 796. The BM receiver performance summary is listed in Table III.

IV. CONCLUSION

We describe a 25 Gb/s dc-coupled BM receiver for optical links in a dynamically reconfigurable network. The architecture of the receiver has two main innovations. The first one is a BM TIA calibration engine, implementing a binary search algorithm to find the input dc current level and settling in 12.5 ns. The second one is a BM-CDR, based on a successive approximation search algorithm, locating edge positions in 18.5 ns. The overall receiver tested with a single mode 1550 nm reference optical transmitter was measured to have $-10.9~\mathrm{dBm}$ sensitivity (average power, BER $<10^{-12}$) at 25 Gb/s, with 4.4 pJ/bit efficiency. The core of the 32 nm SOI CMOS circuit occupies $200~\mathrm{\mu m} \times 300~\mathrm{\mu m}$.

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